

DIACONIS EXERCISE 13

1. PRELIMINARY DEFINITIONS

Let P and Q be probability distributions on a finite set Ω . The *total variation distance* between P and Q , denoted $\|P - Q\|$ is defined by

$$\|P - Q\| = \max_{A \subseteq \Omega} |P(A) - Q(A)|$$

Note that since $P(\Omega \setminus A) - Q(\Omega \setminus A) = -(P(A) - Q(A))$, an equivalent definition is

$$\|P - Q\| = \max_{A \subseteq \Omega} P(A) - Q(A).$$

This definition apparently requires us to consider all events $A \subseteq \Omega$ to find the one on which P and Q assign the most widely differing probabilities. But a moments thought shows that

$$\|P - Q\| = \sum_{\substack{\omega \in \Omega \\ P(\omega) > Q(\omega)}} (P(\omega) - Q(\omega))$$

and hence that

$$(1) \quad \|P - Q\| = \frac{1}{2} \sum_{\omega \in \Omega} |P(\omega) - Q(\omega)|.$$

It still seems quite remarkable to me that (1) can serve as a definition of total variation distance. This equation also shows that total variation distance is essentially the same as the ℓ_1 norm on \mathbf{R}^N .

2. EXPECTATION AND VARIANCE OF THE NUMBER OF FIXED POINTS

Suppose that we shuffle a pack of n cards by choosing uniformly at random two numbers from $\{1, 2, \dots, n\}$. If the numbers are the same, we do nothing; otherwise we swap the cards in the indicated positions. The corresponding probability distribution on the symmetric group S_n is defined by $P(1_{S_n}) = 1/n$ and $P(t) = 2/n^2$ if t is a transposition. If we perform the shuffle k times, then the probability that the cards are permuted according to the permutation $\sigma \in S_n$ is $P^{*k}(\sigma)$. Here P^{*k} is the k -th convolution of P , as defined by $Q^{*1} = Q$ and

$$P^{*k}(\sigma) = \sum_{\pi \in S_n} P^{*(k-1)}(\sigma\pi^{-1})P(\pi)$$

for each $k \geq 2$.

The first part of Exercise 13 in Diaconis' book outlines a proof that if $b > 0$ and $k = \frac{n}{2} \log n - bn$ then the total variation distance between P^{*k} and the uniform distribution U is non-negligible. Earlier, in Theorem 5, it is shown that if $c > 0$ and $k = \frac{n}{2} \log n + cn$ then

$$\|P^{*k} - U\| \leq ae^{-2c}$$

for a constant $a \in \mathbf{R}$, so this result is sharp. Even the crude statement of this 'fast cut-off', that for any $\varepsilon > 0$, a shuffle with $(1/2 + \varepsilon)n \log n$ steps guarantees good mixing, while the shuffle obtained by $(1/2 - \varepsilon)n \log n$ steps will (with overwhelming probability) be poor, seems striking.

Let $F(\sigma)$ denote the number of fixed points of $\sigma \in S_n$. Under the uniform distribution permutations with no fixed points, i.e. derangements, occur with probability about $1/e$. If k is small then it is intuitively clear that the probability distribution P^{*k} will favour permutations with relatively many fixed points, and so derangements, and other permutations with few fixed points, will be underrepresented. By making this precise we shall get a bound in Corollary 9 on the variation distance $\|P^{*k} - U\|$.

We shall in fact solve the more general version of the problem where $P(1_{S_n}) = p_n$ for given $p_n \in \mathbf{R}$, and $P(t) = (1 - p_n) / \binom{n}{2}$ if t is a transposition; this is required to solve the second part of the exercise. For notational convenience, let $q_n = 1 - p_n$.

Proposition 1. *If $k \in \mathbf{N}$ then*

$$\mathbf{E}_{P^{*k}}(F) = 1 + (n - 1) \left(1 - \frac{2q_n}{n - 1}\right)^k.$$

In particular, if $p_n = 1/n$ then $q_n = 1 - 1/n$ and

$$\mathbf{E}_{P^{*k}}(F) = 1 + (n - 1) \left(1 - \frac{2}{n}\right)^k.$$

Proof. The n -dimensional natural representation of S_n decomposes as the sum of an irreducible subrepresentation W of dimension $n - 1$ and the trivial representation. The trace of the permutation matrix representing $\sigma \in S_n$ is simply $F(\sigma)$, so we we have $\text{Tr}_W(\sigma) = F(\sigma) - 1$ for all $\sigma \in S_n$.

Let

$$x = p_n + \frac{q_n}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} (ij) \in \mathbf{CS}_n$$

be the element encoding the probability distribution P . By the previous paragraph we have $\mathbf{E}_P(F - 1) = \text{Tr}_W(x)$. More generally, since the product in the group algebra \mathbf{CS}_n corresponds to the convolution product on probability distributions, we have

$$(2) \quad \mathbf{E}_{P^{*k}}(F - 1) = \text{Tr}_W(x^k)$$

It follows from Lemma 2 below that x acts as $\alpha 1_W$ on W , where

$$\begin{aligned} \alpha &= p_n + \frac{(1-p_n)}{\binom{n}{2}} \binom{n}{2} \frac{\chi_W((12))}{n-1} = p_n + q_n \frac{n-3}{n} \\ &= 1 - \frac{2q_n}{n-1}. \end{aligned}$$

Note that in the particular case where $p_n = 1/n$ we have $\alpha = 1 - \frac{2}{n}$. Hence x^k acts as $(1 - \frac{2}{n-1}(1-p_n))^k 1_W$ on W and, by (2), we have

$$\mathbf{E}_{P^{*k}}(F-1) = (n-1) \left(1 - \frac{2q_n}{n-1}\right)^k$$

from which the result follows immediately. \square

The following lemma, which can be easily proved using Schur's Lemma, was used in the proof of Proposition 1.

Lemma 2. *Let W be an irreducible representation of S_n . If $x \in \mathbf{C}S_n$ is the sum of all elements in the conjugacy class of $\sigma \in S_n$ then x acts on W as*

$$\frac{|x^{S_n}| \chi_W(\sigma)}{\chi_W(1)} 1_W. \quad \square$$

It is easily seen from (1) that if $f : S_n \rightarrow \mathbf{R}$ is any function such that $|f(\sigma)| \leq 1$ for all $\sigma \in S_n$, then

$$\|Q - R\| \geq \frac{1}{2} |f(\sigma)(Q(\sigma) - R(\sigma))|.$$

Taking $f(\sigma) = F(\sigma)/n$ and applying Proposition 1 we get

$$\|P^{*k} - U\| \geq \frac{1}{2n} (\mathbf{E}_{P^{*k}}(F) - 1) = \left(1 - \frac{2q_n}{n-1}\right)^k.$$

In particular, if $p_n = 1/n$ then since $(1 - \frac{2}{n})^n \rightarrow e^{-2}$ as $n \rightarrow \infty$, it follows that n steps do not suffice to get good mixing in this case. The same result holds whenever $p_n \rightarrow 0$ as $n \rightarrow \infty$. As Diaconis remarks, to get a stronger result we need to use the variance of F .

Proposition 3. *If $k \in \mathbf{N}$ and $n \geq 4$ then*

$$\begin{aligned} \mathbf{Var}_{P^{*k}}(F) &= 1 + (n-1) \left(1 - \frac{2q_n}{n-1}\right)^k + \frac{n(n-3)}{2} \left(1 - \frac{4q_n}{n}\right)^k \\ &\quad - \frac{(n-1)(n-2)}{2} \left(1 - \frac{4q_n}{n-1}\right)^k - (n-1)^2 \left(1 - \frac{2q_n}{n-1}\right)^{2k} \end{aligned}$$

Proof. The variance of F is the same as the variance of $F-1$ so as in Proposition 1, we may work with $F-1$. We keep the notation from this proposition. To find the variance of $F-1$ we need the expected value of $(F-1)^2$. Since $(F-1)(\sigma)^2$ is the trace of σ in its action on $W \otimes W$ we have

$$(3) \quad \mathbf{E}_{P^{*k}}(F-1)^2 = \text{Tr}_{W \otimes W}(x).$$

To compute this trace we decompose $W \otimes W$ into its irreducible constituents. We begin by observing that

$$\begin{aligned}\chi_{W \otimes W} &= \chi_W \times \chi_W \\ &= \chi_W \times (1_{S_{n-1}} \uparrow^{S_n} - 1_{S_n}) \\ &= (\chi_W \downarrow_{S_{n-1}}) \uparrow^{S_n} - \chi_W.\end{aligned}$$

In the standard notation for irreducible characters of the symmetric group, $\chi_W = \chi^{(n-1,1)}$. Using the ordinary branching rule (see [1, Chapter 9]) we get

$$\chi_W \downarrow_{S_{n-1}} = \chi^{(n-1)} + \chi^{(n-2,1)}$$

and hence, provided $n \geq 4$,

$$\chi_W \downarrow_{S_{n-1}} \uparrow^{S_n} - \chi_W = \chi^{(n)} + \chi^{(n-1,1)} + \chi^{(n-2,2)} + \chi^{(n-2,1,1)}.$$

Therefore W decomposes as a direct sum of four irreducible representations. By Lemma 2, the scalar by which x acts on an irreducible representation U is

$$p_n + q_n \frac{\chi_U((12))}{\chi_U(1)}.$$

The table below shows $\chi^\lambda(1)$ and $\chi^\lambda(12)$ for the irreducible characters appearing above. These values are easily computed using the Murnaghan–Nakayama rule: see [1, Chapter 21]. The calculations can be simplified by using the identity $(n-3)(n-4)/2 = (n-2)(n-5)/2 + 1$.

λ	$\chi^\lambda(1)$	$\chi^\lambda((12))$
$(n-1, 1)$	$n-1$	$n-3$
$(n-2, 2)$	$n(n-3)/2$	$(n-3)(n-4)/2$
$(n-2, 1, 1)$	$(n-1)(n-2)/2$	$(n-2)(n-5)/2$

It follows that there is a basis of W on which x acts as the matrix

$$I_1 \oplus \left(1 - \frac{2q_n}{n-1}\right) I_{n-1} \oplus \left(1 - \frac{4q_n}{n}\right) I_{n(n-3)/2} \oplus \left(1 - \frac{4q_n}{n-1}\right) I_{(n-1)(n-2)/2}.$$

Hence by (3), we have

$$\begin{aligned}\mathbf{E}_{P^{*k}}(F-1)^2 &= 1 + (n-1) \left(1 - \frac{2q_n}{n-1}\right)^k + \\ &\quad \frac{n(n-3)}{2} \left(1 - \frac{4q_n}{n}\right)^k + \frac{(n-1)(n-2)}{2} \left(1 - \frac{4q_n}{n-1}\right)^k.\end{aligned}$$

The proposition now follows on subtracting

$$\left(\mathbf{E}_{P^{*k}}(F-1)\right)^2 = (n-1)^2 \left(1 - \frac{2q_n}{n-1}\right)^{2k}$$

using the value given in Proposition 1. □

The following special case is worth noting.

Proposition 4. *If $p_n = 1/n$ then*

$$\begin{aligned} \mathbf{Var}_{P^{*k}}(F) = & 1 + (n-1)\left(1 - \frac{2}{n}\right)^k - \frac{(n+1)(n-2)}{2}\left(1 - \frac{2}{n}\right)^{2k} \\ & + \frac{(n-1)(n-2)}{2}\left(1 - \frac{4}{n}\right)^k. \end{aligned}$$

Proof. Substituting $q_n = (n-1)/n$ in Proposition we see that

$$\frac{n(n-3)}{2}\left(1 - \frac{4q_n}{n}\right)^k = \frac{n(n-3)}{2}\left(1 - \frac{2}{n}\right)^{2k}$$

and

$$(n-1)^2\left(1 - \frac{2q_n}{n-1}\right)^{2k} = (n-1)^2\left(1 - \frac{2}{n}\right)^{2k}.$$

The difference of these expressions gives the third term above, and the others come from direct substitution. \square

To get the corollary of Proposition 1 and Proposition 4 when $k = \frac{n}{2} \log n - bn$ we need the following lemma.

Lemma 5. *If $f(n)$ is a polynomial of degree d with leading term an^d and $(d - r_n) \log n \rightarrow 0$ as $n \rightarrow \infty$ then*

$$f(n)\left(1 - \frac{2r_n}{n}\right)^{\frac{n}{2} \log n - bn} = ae^{2bd}\left(1 + O\left(\frac{\log n}{n}\right)\right)$$

as $n \rightarrow \infty$.

Proof. It is not hard to show that $\log f(n) = d \log n + \log a + O(1/n)$ and that

$$\begin{aligned} \log\left(1 - \frac{2r_n}{n}\right)^{\frac{n}{2} \log n - bn} &= \left(\frac{n}{2} \log n - bn\right) \log\left(1 - \frac{2r_n}{n}\right) \\ &= -\left(\frac{n}{2} \log n - bn\right)\left(\frac{2r_n}{n} + O(1/n^2)\right) \\ &= -r_n \log n + 2br_n + O\left(\frac{\log n}{n}\right). \end{aligned}$$

The lemma now follows from the hypothesis that $(d - r_n) \log n \rightarrow 0$ as $n \rightarrow \infty$. \square

Note that the proof of the lemma makes it clear that the implied constant in the $O(1/n)$ term depends on b .

Corollary 6. *Let $k = \frac{n}{2} \log n - bn$. Then*

$$\begin{aligned} \mathbf{E}_{P^{*k}}(F) &\rightarrow 1 + e^{2b}, \\ \mathbf{Var}_{P^{*k}}(F) &\rightarrow 1 + e^{2b}, \end{aligned}$$

provided that $p_n \log n \rightarrow 0$ as $n \rightarrow \infty$,

Proof. By Lemma 5 in the case when $r_n = nq_n/(n-1)$ and $f(n) = (n-1)^2$ we get

$$\lim_{n \rightarrow \infty} (n-1)^2 \left(1 - \frac{2q_n}{n-1}\right)^k = e^{2b}.$$

To apply the lemma we need that

$$\left(1 - \frac{nq_n}{n-1}\right) \log n \rightarrow 0 \quad \text{as } n \rightarrow \infty;$$

this holds because

$$-\frac{\log n}{n} \leq \left(1 - \frac{nq_n}{n-1}\right) \log n \leq p_n \log n$$

and by assumption $p_n \log n \rightarrow 0$ as $n \rightarrow \infty$. Hence, by Proposition 1 we have

$$\lim_{n \rightarrow \infty} \mathbf{E}_{P^{*k}}(F) = 1 + e^{2b}.$$

The proof of the limit for $\mathbf{Var}_{P^{*k}}(F)$ is similar using the lemma and Proposition 2. \square

3. BOUNDS

The following two propositions will be used to turn this corollary into a bound on $\|P^{*k} - U\|$.

Proposition 7. *Suppose that $p_n \log n \rightarrow 0$ as $n \rightarrow \infty$. Let $M \in \mathbf{N}$. If $k = \frac{n}{2} \log n - bn$ where $b \geq 2$ and $M = \lfloor e^{2b}/2 \rfloor$ then*

$$\mathbf{P}_{P^{*k}}(F \leq M) \leq \frac{2}{M}$$

for all k sufficiently large.

Proof. Choose k sufficiently large so that $\mathbf{E}_{P^{*k}}(F) \geq e^{2b}$ and $\mathbf{Var}_{P^{*k}}(F) \leq 2 + e^{2b}$. Since $(1 + e^b)^2 \geq 2 + e^{2b}$, it follows from Chebychev's inequality that

$$\mathbf{P}_{P^{*k}}\left(F \leq e^{2b} - t(1 + e^b)\right) \leq \frac{1}{t^2}.$$

Putting $t = \frac{1}{2}(e^b - 1)$ we get

$$\mathbf{P}_{P^{*k}}\left(F \leq \frac{e^{2b}}{2}\right) \leq \frac{2}{(e^b - 1)^2}.$$

One can check that

$$(e^b - 1)^2 \geq e^{2b}/2$$

for all $b \geq 2$. It follows that if $M = \lfloor e^{2b}/2 \rfloor$ then

$$\mathbf{P}_{P^{*k}}(F \leq M) \leq \mathbf{P}_{P^{*k}}\left(F \leq \frac{e^{2b}}{2}\right) \leq \frac{2}{(e^b - 1)^2} \leq \frac{2}{e^{2b}/2} \leq \frac{2}{M}$$

as required. \square

Proposition 8. *Let $M \in \mathbf{N}$. Then*

$$\mathbf{P}_U(F > M) \leq \frac{1}{(M+1)!}.$$

Proof. The probability that a particular $M+1$ -subset of $\{1, 2, \dots, n\}$ is fixed by a permutation in S_n is $(n-M-1)!/n!$; now sum over all $M+1$ -subsets to get

$$\mathbf{P}_U(F > M) \leq \binom{n}{M+1} \frac{(n-M-1)!}{n!} = \frac{1}{(M+1)!}. \quad \square$$

Corollary 9. *Suppose that $p_n \log n \rightarrow 0$ as $n \rightarrow \infty$. If $k = \frac{n}{2} \log n - bn$ where $b \geq 2$ then*

$$\|P^{*k} - U\| \geq 1 - \frac{6}{e^{2b} - 2}$$

provided k is sufficiently large.

Proof. Let $M = \lfloor e^{2b}/2 \rfloor$ and consider the event $F \leq M$. We have

$$\|P^{*k} - U\| \geq P_U(F \leq M) - P_{P^{*k}}(F \leq M).$$

It follows from the previous two propositions that, provided k is sufficiently large

$$\|P^{*k} - U\| \geq 1 - \frac{1}{(M+1)!} - \frac{2}{M} \geq 1 - \frac{3}{M} \geq 1 - \frac{6}{e^{2b} - 2}$$

as required. □

This gives a non-trivial bound provided that $b \geq 1.04$. Therefore provided $p_n = O(1/\log n)$ we have shown that $\frac{n}{2} \log n - 2n$ steps do not suffice to get good mixing.

4. SECOND PART OF EXERCISE 13

The second part of Exercise 13 claims that if $p_n = 1/(1 + \binom{n}{2})$, so the identity is equally likely to be chosen as any transpositions, then $c(n)n^2$ are necessary to get a good shuffle, where $c(n) \rightarrow \infty$ as $n \rightarrow \infty$.

It follows easily from Proposition 1 and the inequality $1 - x \leq e^{-x}$ that

$$0 \leq \mathbf{E}_{P^{*c(n)n^2}}(F - 1) \leq (n - 1)e^{-2c(n)nq_n}$$

for all $n \in \mathbf{N}$. A similar bound will hold for the variance of F . So it seems that if $k = c(n)n^2$ then F is not able to detect any significant difference between P^{*k} and U unless p_n is at least $1/\log n$; certainly the smaller value $p_n = 1/(1 + \binom{n}{2})$, as compared to $1/n$, will not cause any unexpected problems.

It is however possible to use similar ideas to get a slightly weaker result. The key observation is that if we make much fewer than $1/p_n$ steps then there is a good chance that we have never chosen the identity. In this case,

the sign of the resulting permutation is given the parity of the number of steps. So if $A_n \subseteq S_n$ is the alternating group then

$$\mathbf{P}_{P^{*k}}(A_n) \geq (1 - p_n)^k \geq 1 - kp_n$$

whenever k is even. Since $\mathbf{P}_U(A_n) = 1/2$ for all n , it follows from the definition of total variation distance that

$$\|P^{*k} - U\| \geq 1/2 - kp_n$$

whenever k is even. So in this particular case we have

$$\|P^{*k} - U\| \geq 1/2 - \theta$$

whenever k is even and $k \leq \theta n^2/2$. So $n^2/4$ steps do not suffice.

REFERENCES

- [1] JAMES, G. D. *The representation theory of the symmetric groups*, vol. 682 of *Lecture Notes in Mathematics*. Springer, Berlin, 1978.